

Claim Amendments:

1. (currently amended) A method comprising the steps of:

encrypting a data message  $m$  at a transmitter processor using a primary transmitter secret key  $z$ , wherein  $z$  is known to the transmitter processor but not to a receiver processor, to form a quantity  $E$ , wherein El-Gamal encryption is used for encrypting the data message  $m$ ;

preparing a quadruplet  $(a_{\text{new}}, b_{\text{new}}, s_{\text{new}}, E)$  at the transmitter processor where:

$$a_{\text{new}} = z^* y^c \text{ modulo } p;$$

$$b_{\text{new}} = g^c \text{ modulo } p;$$

$$s_{\text{new}} = \text{signature}_c(a_{\text{new}}, b_{\text{new}}, E);$$

where  $y = g^x \text{ modulo } p$ ,  $c$  is a random number which is used in the step of encrypting the data message  $m$  using El-Gamal encryption,  $x$  is a receiver secret key of the receiver processor, and the parameters  $g$ ,  $x$ , and  $p$  are picked using a known encryption method;

wherein  $s_{\text{new}}$  is a signature which is determined by using the same random number  $c$  that was used to determine  $a_{\text{new}}$  and  $b_{\text{new}}$ ;

transmitting the quadruplet  $(a_{\text{new}}, b_{\text{new}}, s_{\text{new}}, E)$  from the transmitter processor to the receiver processor;

verifying the signature  $s_{\text{new}}$  at the receiver processor;

decrypting  $a_{\text{new}}$  and  $b_{\text{new}}$  at the receiver processor by using the receiver secret key  $x$  to get the primary transmitter secret key  $z$ ;

using the primary transmitter secret key  $z$  to decrypt the quantity  $E$  and thereby obtaining the message  $m$  at the receiver processor.

2. (currently amended) The method of claim 1 and wherein:

the step of decrypting  $a_{\text{new}}$  and  $b_{\text{new}}$  at the receiver processor using the receiver secret key  $x$  to get the primary transmitter secret key  $z$  is comprised of computing  $z = a_{\text{new}}/b_{\text{new}}^x$ .

3. (cancelled)

4. (cancelled)

5. (currently amended) The method of claim 1 wherein:

the primary transmitter secret key  $z$  is determined at the transmitter processor from the formula of  $z = g^Y$  modulo  $p$ , where  $Y$  is a random value chosen from the set  $[0..q]$ , where  $q$  is a value picked using a known encryption method.

6. (currently amended) A method comprising the steps of:

creating a primary transmitter key  $z$  at a transmitter processor wherein the primary transmitter key is known to the transmitter processor but not to a receiver processor;

creating a secondary transmitter key  $z'$  at the transmitter processor wherein the secondary transmitter key is known to the transmitter processor but not to the receiver processor, wherein the secondary transmitter key  $z'$  which is a function of  $z$ ;

encrypting a data message  $m$  at the transmitter processor, using the secondary transmitter secret key  $z'$  to form a quantity  $E$  wherein El Gamal encryption is used for encrypting the data message  $m$ ;

preparing a quadruplet  $(a_{new}, b_{new}, s_{new}, E)$  at the transmitter processor, where:

$$a_{new} = z^* y^c \text{ modulo } p;$$

$$b_{new} = g^c \text{ modulo } p;$$

$$s_{new} = \text{signature}_c(a_{new}, b_{new}, E);$$

where  $y = g^x$  modulo  $p$ ,  $c$  is a random number which is used in the step of encrypting the data message  $m$  using El Gamal encryption,  $x$  is a receiver secret key of the receiver processor, and the parameters  $g$ ,  $x$ , and  $p$  are picked using a known encryption method;

wherein  $s_{new}$  is a signature which is determined by using the same random number  $c$

that was used to determine  $a_{new}$  and  $b_{new}$ :

transmitting the quadruplet ( $a_{new}$ ,  $b_{new}$ ,  $s_{new}$ ,  $E$ ) from the transmitter processor to the receiver processor;

verifying the signature  $s_{new}$  at the receiver processor;

decrypting  $a_{new}$  and  $b_{new}$  at the receiver processor, using the receiver secret key  $x$  to get the primary transmitter secret key  $z$ ;

modifying the primary transmitter secret key  $z$  at the receiver processor, to obtain the secondary transmitter secret key  $z'$  and using the secondary transmitter secret key  $z'$  to decrypt the quantity  $E$  and thereby obtaining the message  $m$  at the receiver processor.

7. (original) The method of claim 6 and wherein:

the primary transmitter key  $z$  is provided which is not of the format used for producing the ciphertext  $E$ ;

the secondary transmitter key  $z'$  is computed as a function of  $z$ , where the function is an arbitrary function.

8. (currently amended) A method comprising the steps of:

creating a primary transmitter key  $z$  at a transmitter processor;

creating a secondary transmitter key  $z'$  which is a function of  $z$  at the transmitter processor;

providing a plurality of portion keys which are derived from the secondary transmitter key  $z'$  at the transmitter processor;

encrypting a data message  $m$  at the transmitter processor, using the plurality of portion keys to form a quantity  $E$  wherein ElGamal encryption is used for encrypting the data message  $m$ ;

preparing a quadruplet ( $a_{new}$ ,  $b_{new}$ ,  $s_{new}$ ,  $E$ ) at the transmitter processor, where:

$$a_{new} = z^* y^c \text{ modulo } p;$$

$$b_{new} = g^c \text{ modulo } p;$$

$$s_{new} = \text{signature}_c(a_{new}, b_{new}, E);$$

where  $y = g^x \text{ modulo } p$ ,  $c$  is a random number ~~which is used in the step of encrypting the data message  $m$  using El-Gamal encryption~~,  $x$  is a receiver secret key ~~of a receiver processor~~, and the parameters  $g$ ,  $x$ , and  $p$  are picked using a known encryption method;

wherein  $s_{new}$  is a signature which is determined by using the same random number  $c$  that was used to determine  $a_{new}$  and  $b_{new}$ :

transmitting the quadruplet ( $a_{new}$ ,  $b_{new}$ ,  $s_{new}$ ,  $E$ ) from the transmitter processor to the receiver processor;

verifying the signature  $s_{new}$  at a the receiver processor;

decrypting  $a_{new}$  and  $b_{new}$  at the receiver processor, using the receiver secret key  $x$  to get the primary transmitter secret key  $z$ ;

modifying the primary transmitter secret key  $z$  at the receiver processor, to obtain the secondary transmitter secret key  $z'$  and using the secondary transmitter secret key  $z'$  to determine the plurality of portion keys and using the plurality of portion keys to decrypt the quantity  $E$  and thereby obtaining the message  $m$  at the receiver processor.

9. (previously presented) The method of claim 1 wherein

the signature  $s_{new}$  is determined by using a Schnorr signature method.

10. (previously presented) The method of claim 1 wherein

the signature  $s_{new}$  is determined using a Digital Signature Standard.

11. (currently amended) An apparatus comprising

a transmitter processor;

wherein the transmitter processor

encrypts a data message  $m$  using a primary transmitter secret key  $z$ ,

known to the transmitter processor but not known to a receiver processor, to

form a quantity  $E$ , wherein El-Gamal encryption is used to encrypt the data

message  $m$ ; and

prepares a quadruplet  $(a_{\text{new}}, b_{\text{new}}, s_{\text{new}}, E)$  where:

$$a_{\text{new}} = z^x y^c \text{ modulo } p;$$

$$b_{\text{new}} = g^c \text{ modulo } p;$$

$$s_{\text{new}} = \text{signature}_c(a_{\text{new}}, b_{\text{new}}, E);$$

where  $y = g^x \text{ modulo } p$ ,  $c$  is a random number ~~which is used in the step of~~

~~encrypting the data message  $m$  using El-Gamal encryption~~,  $x$  is a receiver secret key of

~~the receiver processor~~, and the parameters  $g$ ,  $x$ , and  $p$  are picked using a known

encryption method; and

wherein  $s_{\text{new}}$  is a signature, and wherein the transmitter processor determines

$s_{\text{new}}$  by using the same random number  $c$  that was used to determine  $a_{\text{new}}$  and  $b_{\text{new}}$ .

12. (cancelled).

13. (currently amended) The apparatus of claim 11 wherein

the transmitter processor uses a Schnorr signature method to determine  $s_{\text{new}}$ .

14. (currently amended) The apparatus of claim 11 wherein

the transmitter processor uses a Digital Signature Standard to determine  $s_{\text{new}}$ .

15. (new) The method of claim 1 wherein

El Gamal encryption is used for the encrypting steps.

16. (new) The method of claim 6 wherein

El Gamal encryption is used for the encrypting steps.

17. (new) The method of claim 8 wherein

El Gamal encryption is used for the encrypting steps.

18. (new) The apparatus of claim 11 wherein

El Gamal encryption is used for encrypting.

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